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DiCAP - An Architecture for Distributed Packet Capturing

Cristian Morariu

Department of Informatics IFI, Communication Systems Group CSG, University of Zürich



Motivation
Design and Implementation
Evaluation
Concluding Remarks



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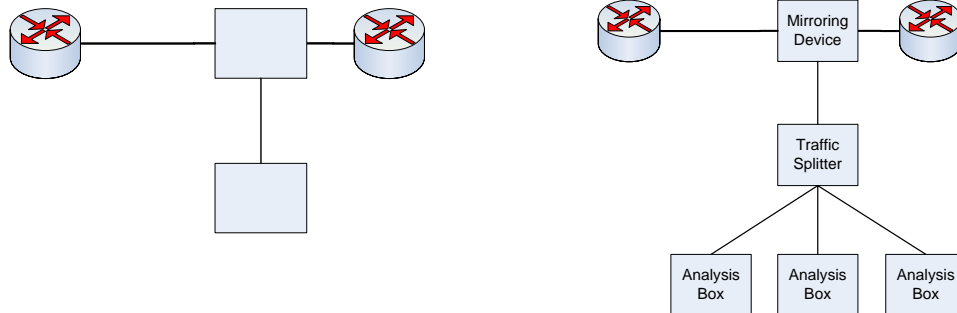
Overview

- ❑ Most network monitoring tasks require packet inspection
- ❑ Network monitoring
 - *Live*
 - packets are inspected in real-time
 - data is dropped after inspection
 - e.g. traffic accounting, IDS
- ❑ High packet rates reduce the time available to handle a single packet
- ❑ Solutions for high packet rates:
 - packet sampling
 - decreased measurement accuracy
 - dedicated hardware
 - Expensive

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Traditional Architecture for Traffic Analysis



- ❑ Simple
- ❑ Easy to deploy
- ❑ Not scalable due to a single analysis box
- ❑ More complex
- ❑ More scalable with respect to performance
- ❑ The single point of failure still present

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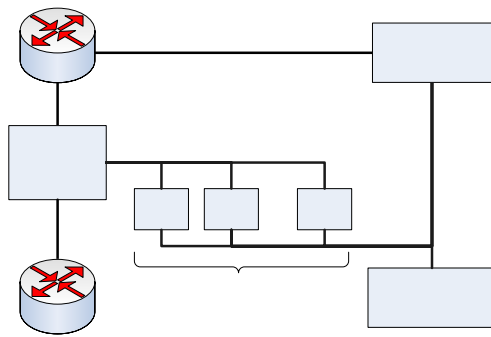
Motivation

- ❑ Build a **distributed** architecture for IP traffic **capturing**
 - Avoid a single point of failure
 - Avoid dedicated hardware
 - Based on stock, inexpensive PCs
 - Allows the increase of packets that can be captured

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DiCAP Architecture



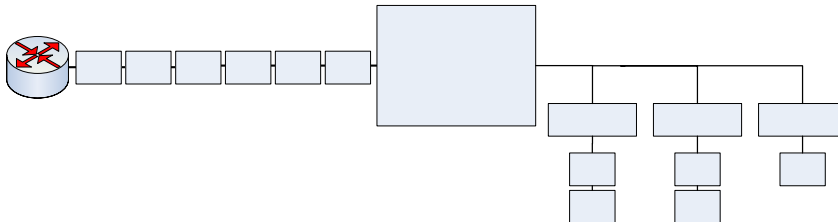
- ❑ Each node receives all packets
- ❑ Node coordinator decides the capture policy on each capture node
- ❑ The capture nodes select which packets to capture so that no packet is captured by two nodes
- ❑ Packet data is sent to the packet data analyzer for further analysis

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Selection Policies (1)

- ❑ Round robin selection
 - Node coordinator introduces control packets in the mirrored traffic
 - Capture nodes are logically organized in a chain
 - After each packet, capture responsibility shifts to the next node in the chain
 - Node coordinator configures the chain



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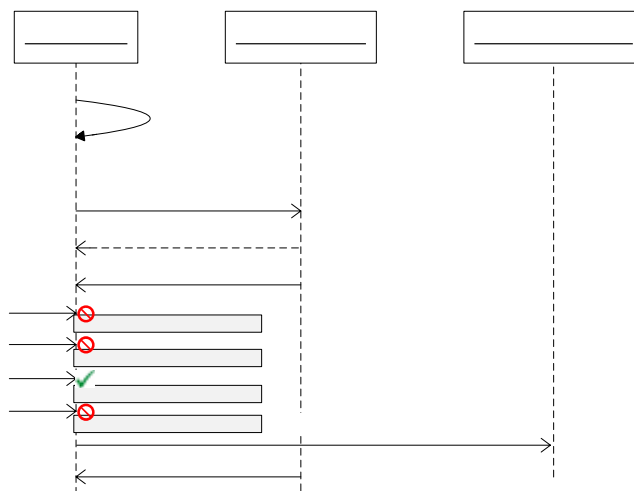
Selection Policies (2)

- Hash-based selection
 - A hash function is applied on packet headers
 - Current implementation uses IP identification field
 - Each node is responsible with a particular range of hash values
 - Each packet is captured by the node responsible with the respective range of hash values
- Advantages:
 - Easier synchronization
 - No need of control packet injection
- Disadvantages:
 - More computation needed for hash calculation

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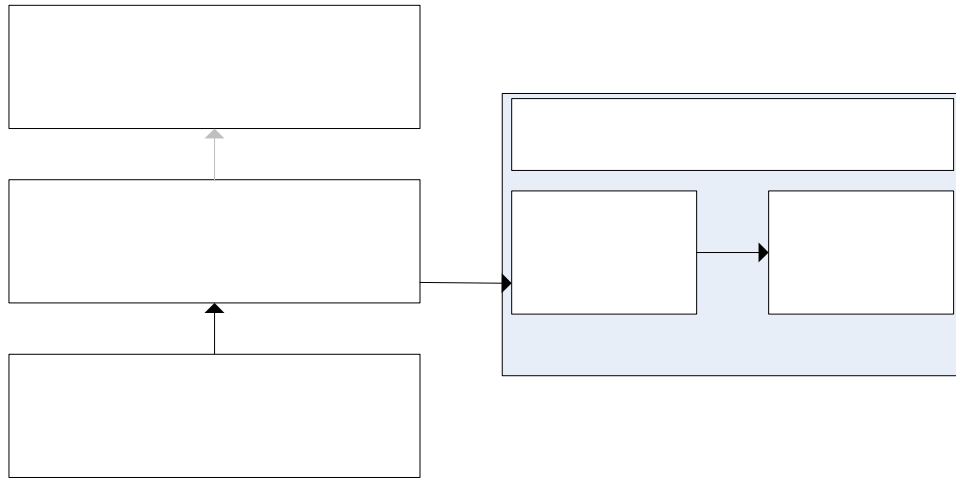
Capture Example in Round-Robin Selection



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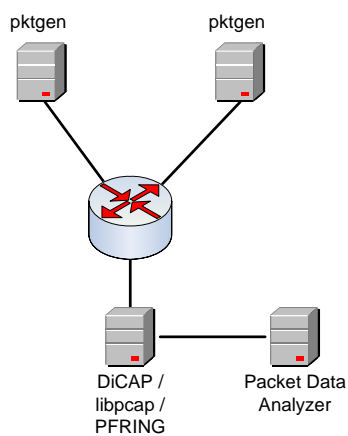
DiCAP Implementation



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Evaluation Testbed



- Two nodes used for traffic injection using Linux pktgen
- One DiCAP node for testing capture performance of a single node
 - 1G Ethernet Broadcom NetXtreme BCM5721 card
- Packet data is sent to the packet data analyzer for further analysis
- One packet data analyzer

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Linux

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DiCAP Capture Evaluation

Packet Rate	libpcap Loss	PFRING Loss	DiCAP Loss
119 Kpps	0%	0%	0%
232 Kpps	10%	1%	0%
380 Kpps	75%	28%	0%
492 Kpps	90%	83%	0%
620 Kpps	93%	96%	0%

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Operation Modes

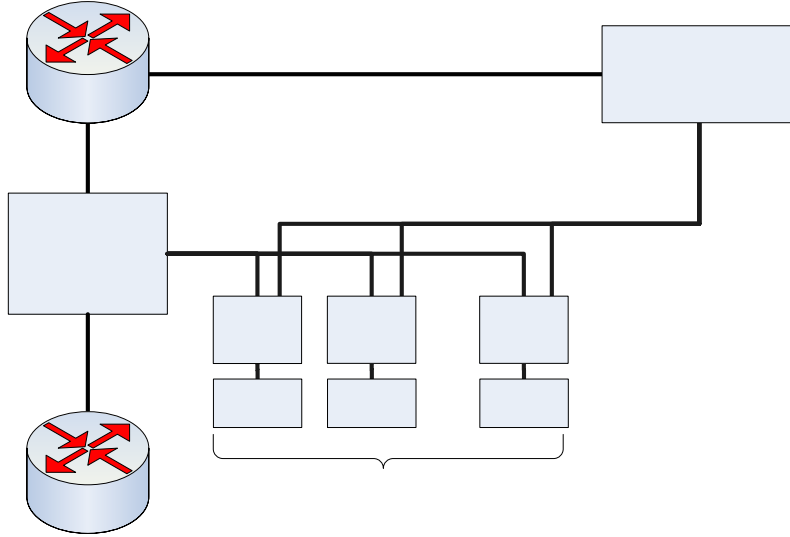
- Capture Mode
 - Packet headers are captured and forwarded to a packet analyzer
 - No possibility to analyze payload of packets

- Distribution Mode
 - DiCAP does not capture anything
 - Decides which packets are forwarded to the higher levels
 - Capture takes place in user space applications
 - Allows parallel use of *libpcap*

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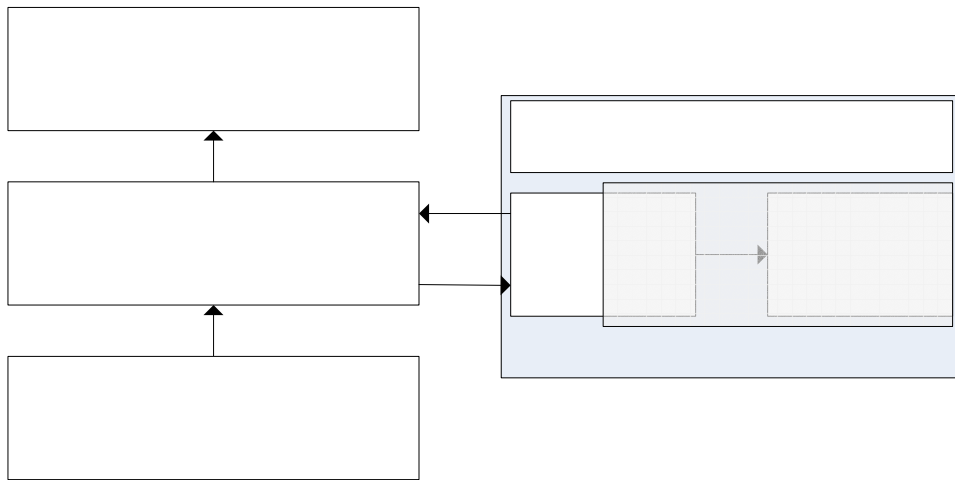
DiCAP in Distribution Mode



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DiCAP Implementation – Distribution Mode

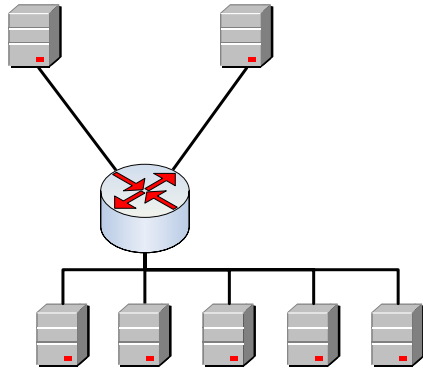


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Mirrorin
Device

Evaluation Testbed

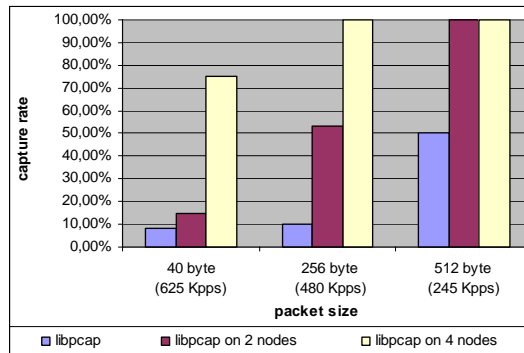


- Two nodes used for traffic injection using Linux pktgen
- One libpcap node for testing
- Up to 4 DiCAP nodes in distribution mode

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Distribution Mode Evaluation



- 3 different tests were performed
 1. traditional libpcap on a single PC
 2. DiCAP in distribution mode on 2 PCs
 3. DiCAP in distribution mode on 4 PCs
- DiCAP can improve the capture performance up to 700% when 4 PCs are used in parallel

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Concluding Remarks

- ❑ DiCAP can be used to distributedly capture packets on a high-speed link
- ❑ It may be used to allow a distributed deployment of libpcap-based applications running on common, inexpensive PCs
- ❑ It may significantly improve libpcap performance
- ❑ Very simple design, easy to implement in hardware.
- ❑ Further investigation on using other hash functions may lead to performance improvement and better load ballance

Thank You!
